

IN THE CLAIMS:

1. (Original) Method for operating (implementing) a secondary operating system on a computer in addition to a primary operating system, wherein a secondary operating system driver (SOS driver) of the primary operating system is loaded for loading and controlling the secondary operating system.

2. (Original) Method according to claim 1, wherein the secondary operating system driver (SOS driver) subsequently loads the secondary operating system.

3. (Currently Amended) Method according to claim 1 ~~[[or 2]]~~, wherein the secondary operating system driver loads the secondary operating system.

4. (Currently Amended) Method according to ~~one of the claim~~[[s]] 1 ~~[[to 3]]~~, wherein memory contexts (virtual operating areas) are created in the central processing unit (CPU) .

5. (Currently Amended) Method according to ~~one of the claim~~[[s]] 1 ~~[[to 4]]~~, wherein a change between the operating systems takes place by means of the SOS driver of the primary operating system and the board support package (BSP) .

6. (Currently Amended) Method, particularly according to ~~one of the preceding~~ claim[[s]] 1, wherein there is an exchange of interrupt tables on changing the dependence of the

operating systems.

7. (Currently Amended) Method according to ~~one of the~~ claim[[s]] 1 [[to 6]], wherein the secondary operating system controls a change to the primary operating system.

8. (Original) Method according to claim 7, wherein a change from the secondary operating system to the primary operating system takes place when the secondary operating system is idle (entry into the idle loop).

9. (Original) Method according to claim 8, wherein a change from the secondary operating system to the primary operating system takes place through an instruction in the program sequence of the secondary operating system.

10. (Currently Amended) Method according to ~~one of the~~ claim[[s]] 1 [[to 9]], wherein a change from the primary operating system to the secondary operating system takes place through an interrupt call.

11. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein the change between operating systems takes place by means of a program code filed in the tunnel area of the memory.

12. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein interrupt calls of the primary operating system are inhibited during the secondary operating system sequence.

13. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein an interrupt servicing routine in the SOS operator reads the interrupt call table of the secondary operating system and the processing of the latter takes place or is continued at the point relative to the interrupt call.

14. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein for each interrupt associated with the secondary operating system (i.e. which is to initiate an interrupt call in the secondary operating system) the system driver generates an entry in the interrupt call table in the primary operating system, which in turn initiates a call of the
5 corresponding interrupt servicing routine in the secondary operating system.

15. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein by means of an interrupt call servicing routine in the system driver, the information stored in the interrupt table of the secondary operating system (SOS) is determined as to the point in the latter where the running of the interrupt is to take place.

16. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1,

wherein in the case of activity of the secondary operating system (SOS) following an interrupt request through the information stored in the interrupt call table of the secondary operating system as to the point in the latter where the running of the interrupt is to take place, the interrupt call servicing routine of the secondary operating system (SOS) is directly polled solely by the secondary operating system and not via the system driver.

17. (Currently Amended) Method according to ~~one of the claim~~[[s]] 12 [[to 16]], wherein after occurrence a corresponding interrupt call and determination of the point in the secondary operating system where interrupt running is to take place is determined, processing thereof at the point in the secondary operating system concerning the interrupt call is continued.

18. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein on changing from one operating system to the other all the system states of one operating system are stored.

19. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein on changing from one operating system to the other all system states of the other operating system are loaded.

20. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1,

wherein clock generation for the secondary operating system takes place through the hardware timer.

21. (Currently Amended) Method according to ~~one of the preceding~~ claim[[s]] 1, wherein clock generation for the primary operating system takes place through a clock system driver.

22. (Original) Device for operating a secondary operating system on a computer in addition to a primary operating system, wherein a secondary operating system driver (SOS driver) of the primary operating system for loading and controlling the secondary operating system is constructed.

23. (Original) Device according to claim 22, wherein the SOS driver has a tunnel context setting routine for setting a tunnel context in the central processing unit (CPU).

24. (Currently Amended) Device, particularly according to claim 22 [[or 23]], wherein said device is constructed for exchanging interrupt tables on a change of activity of the operating systems.

25. (Currently Amended) Device according to claim 22 [[or 24]], wherein the SOS driver has an interrupt call table change routine for producing entries in the interrupt call table

of the primary operating system, which at least take up entries for the interrupt calls for the secondary operating system.

26. (Currently Amended) Device according to ~~one of the~~ claim[[s]] 22 [[to 25]], wherein the board support package (BSP) has a section for return to the primary operating system (POS).

27. (Currently Amended) Device according to ~~one of the~~ claim[[s]] 22 [[to 26]], wherein the secondary operating system driver (SOS driver) has an interrupt table section by means of which it produces in the primary operating system an interrupt call table containing a call of an interrupt servicing routine for polling the secondary operating system.

28. (Currently Amended) Device according to ~~one of the~~ claim[[s]] 22 [[to 27]], wherein the system driver is constructed for producing an entry in the interrupt call table in the primary operating system (POS) for each interrupt associated with the secondary operating system (SOS), which i.e. is intended to initiate an interrupt call in the secondary operating system and that the interrupt call table is constructed for polling the corresponding interrupt servicing routine in the secondary operating system (SOS).

29. (Currently Amended) Device according to ~~one of the~~ claim[[s]] 22 [[to 28]], wherein an interrupt call servicing routine in the system driver is constructed for determining

the information stored in the SOS interrupt table as to the point in the secondary operating system where interrupt running is to take place.

30. (Currently Amended) Device according to ~~one of the claim~~[[s]] 22 [[to 29]], wherein it is constructed in the case of activity of the secondary operating system (SOS) following an interrupt request through the information stored in the secondary operating system interrupt call table as to the point in which the secondary operating system interrupt running is to take place, so as to poll the interrupt call servicing routine of the secondary operating system directly solely through the secondary operating system and without passing via the system driver.

31. (New) Method for operating (implementing) a secondary operating system on a computer in addition to a primary operating system in which a change from the primary operating system to the secondary operating system takes place through an interrupt call, wherein a secondary operating system driver (SOS driver) in the primary operating system is loaded and activated for loading and controlling the secondary operating system, wherein by means of an interrupt call servicing routine in the system driver said secondary operating system driver (SOS driver) determines the information stored in the interrupt table of the secondary operating system (SOS) as to the point in the latter where the running of the interrupt is to take place.

32. (New) Device for operating a secondary operating system on a computer in addition to a primary operating system in which a change from the primary operating system to the secondary operating system takes place through an interrupt call, wherein it is constructed by a secondary operating system driver (SOS driver) of the primary operating system for loading and controlling the secondary operating system, wherein for execution of the interrupt routine an interrupt call servicing routine in the system driver is constructed for determining the information stored in the interrupt table of the secondary operating system as to the point in the latter where interrupt running is to take place.